



## Runge módszer

**Tengely Szabolcs**

2011. április 15.

Habilitációs előadás, Debrecen





---

## Runge feltétel

Legyen

$$P(X, Y) = \sum_{i=0}^m \sum_{j=0}^n a_{i,j} X^i Y^j,$$

ahol  $a_{i,j} \in \mathbb{Z}$  és  $m > 0, n > 0$ , irreducibilis polinom a  $\mathbb{Q}[X, Y]$  gyűrűben.



---

## Runge feltétel

Legyen

$$P(X, Y) = \sum_{i=0}^m \sum_{j=0}^n a_{i,j} X^i Y^j,$$

ahol  $a_{i,j} \in \mathbb{Z}$  és  $m > 0, n > 0$ , irreducibilis polinom a  $\mathbb{Q}[X, Y]$  gyűrűben.

1887-ben Runge diofantikus egyenletek egy osztályára adott effektív végességi tételt.

$$P(X, Y) = 0.$$





# Runge feltétel

## Runge feltétel

$P$  kielégíti a Runge feltételt, ha a következő kritériumok közül legalább az egyik nem teljesül:

- $X$  és  $Y$  legmagasabb kitevős hatványai izoláltan fordulnak elő,  $aX^m$  és  $bY^n$  alakban,
- minden  $a_{i,j}X^iY^j$  esetén  $ni + mj \leq mn$ ,
- $\sum_{ni+mj=mn} a_{i,j}X^iY^j$  egy irreducibilis polinom hatványa konstans szorzótól eltekintve.





---

## Runge feltétel

Legyen  $\lambda > 0$ .

- $P_\lambda(X, Y)$  : az összes  $a_{i,j}X^iY^j$  összege, ahol  $i + \lambda j$  maximális,
- $\tilde{P}(X, Y)$  : az összes olyan  $a_{i,j}X^iY^j$  összege, amely előfordul valamely  $P_\lambda(X, Y)$  polinomban.

Egy  $P$  polinom eleget tesz Runge feltételének, ha nem létezik  $\lambda$ , amelyre  $\tilde{P} = P_\lambda$  konstans szorzótól eltekintve egy irreducibilis polinom hatványa.



# Runge feltétel

Legyen  $\lambda > 0$ .

- $P_\lambda(X, Y)$  : az összes  $a_{i,j}X^iY^j$  összege, ahol  $i + \lambda j$  maximális,
- $\tilde{P}(X, Y)$  : az összes olyan  $a_{i,j}X^iY^j$  összege, amely előfordul valamely  $P_\lambda(X, Y)$  polinomban.

Egy  $P$  polinom eleget tesz Runge feltételének, ha nem létezik  $\lambda$ , amelyre  $\tilde{P} = P_\lambda$  konstans szorzótól eltekintve egy irreducibilis polinom hatványa.

## Tétel (Runge (1887))

Ha a  $P$  polinom eleget tesz a Runge feltételnek, akkor a  $P(x, y) = 0$  diofantikus egyenletnek csak véges sok megoldása van.





---

## Példa

$$P(X, Y) = X^2 - Y^8 - Y^7 - Y^2 - 3Y + 5,$$

- $P_\lambda(X, Y) = X^2$  ha  $\lambda < \frac{1}{4}$ ,
- $P_\lambda(X, Y) = X^2 - Y^8$  ha  $\lambda = \frac{1}{4}$ ,
- $P_\lambda(X, Y) = Y^8$  ha  $\lambda > \frac{1}{4}$ ,

így  $\tilde{P}(X, Y) = X^2 - Y^8 = (X - Y^4)(X + Y^4)$ .





## Kapcsolódó eredmények

Schinzel (1969)

Masser (1980)

Hilliker és Straus (1983)

Ayad (1991)

Grytczuk és Schinzel (1991)

Walsh (1992)

Poulakis (1999)

Szalay (2000, 2002)

Tengely (2003)

Laurent és Poulakis (2004)

Levin (2004)

Beukers és Tengely (2005)

Sankaranarayanan és Saradha (2008)





# Korlát a megoldásokra

$F(x) = G(y)$ ,  $\deg F = n$ ,  $\deg G = m$ ,  $\gcd(m, n) > 1$  és az  $F(X) - G(Y)$  polinom irreducibilis a  $\mathbb{Q}[X, Y]$  gyűrűben. Legyen  $1 < d \mid \gcd(m, n)$ . Ekkor teljesül Runge feltétele a  $P(X, Y) = F(X) - G(Y)$  polinomra.

## Tétel (T.Sz.)

Ha  $(x, y) \in \mathbb{Z}^2$  megoldása az  $F(x) = G(y)$  egyenletnek, ahol  $F$  és  $G$  kielégítik a fenti feltételeket, akkor

$$\max\{|x|, |y|\} \leq d^{\frac{2m^2}{d}-m}(m+1)^{\frac{3m}{2d}}\left(\frac{m}{d}+1\right)^{\frac{3m}{2}}(h+1)^{\frac{m^2+mn+m}{d}+2m}.$$





# A bizonyítás lépései

## Lemma (Walsh)

Az  $U^d = F(X)$ ,  $V^d = G(X)$  egyenletek által definiált  $U, V$  algebrai függvényekhez léteznek Puiseux kifejtések:

$$u(X) = \sum_{i=-\frac{n}{d}}^{\infty} f_i X^{-i} \text{ és} \quad v(X) = \sum_{i=-\frac{m}{d}}^{\infty} g_i X^{-i}$$

úgy, hogy  $d^{2(n/d+i)-1} f_i \in \mathbb{Z}$  minden  $i > -\frac{n}{d}$  esetén, hasonlóan  $d^{2(m/d+i)-1} g_i \in \mathbb{Z}$  minden  $i > -\frac{m}{d}$  esetén, és  $f_{-\frac{n}{d}} = g_{-\frac{m}{d}} = 1$ . Továbbá  $|f_i| \leq (H(F) + 1)^{\frac{n}{d}+i+1}$  ha  $i \geq -\frac{n}{d}$  és  $|g_i| \leq (H(G) + 1)^{\frac{m}{d}+i+1}$  ha  $i \geq -\frac{m}{d}$ .





# A bizonyítás lépései

$$F(X) = \left( \sum_{i=-\frac{n}{d}}^{\infty} f_i X^{-i} \right)^d, \quad G(Y) = \left( \sum_{i=-\frac{m}{d}}^{\infty} g_i Y^{-i} \right)^d$$

Amennyiben  $|t|$  elég nagy:

$$\left| \sum_{i=1}^{\infty} d^{\frac{2m}{d}-1} f_i t^{-i} \right| < \frac{1}{2}$$

és

$$\left| \sum_{i=1}^{\infty} d^{\frac{2m}{d}-1} g_i t^{-i} \right| < \frac{1}{2}.$$





---

## A bizonyítás lépései

Mivel  $F(x) = G(y)$ , így  $u(x)^d - v(y)^d = 0$ , azaz

$$(u(x) - v(y)) (u(x)^{d-1} + u(x)^{d-2}v(y) + \dots + v(y)^{d-1}) = 0,$$

ha  $d$  páratlan,

$$(u(x)^2 - v(y)^2) (u(x)^{d-2} + u(x)^{d-4}v(y)^2 + \dots + v(y)^{d-2}) = 0,$$

ha  $d$  páros.



---

## A bizonyítás lépései

Mivel  $F(x) = G(y)$ , így  $u(x)^d - v(y)^d = 0$ , azaz

$$(u(x) - v(y)) (u(x)^{d-1} + u(x)^{d-2}v(y) + \dots + v(y)^{d-1}) = 0,$$

ha  $d$  páratlan,

$$(u(x)^2 - v(y)^2) (u(x)^{d-2} + u(x)^{d-4}v(y)^2 + \dots + v(y)^{d-2}) = 0,$$

ha  $d$  páros.

A fenti egyenletekből adódik, hogy

$$u(x) = v(y) \text{ ha } d \text{ páratlan, és}$$

$$u(x) = \pm v(y) \text{ ha } d \text{ páros.}$$





## A bizonyítás lépései

Így

$$0 = |u(x) \pm v(y)| = \left| \sum_{i=-\frac{n}{d}}^{\infty} f_i x^{-i} \pm \sum_{i=-\frac{m}{d}}^{\infty} g_i y^{-i} \right|.$$

Elég nagy  $|x|$  és  $|y|$  esetén

$$\left| \sum_{i=-\frac{n}{d}}^0 d^{\frac{2m}{d}-1} f_i x^{-i} \pm \sum_{i=-\frac{m}{d}}^0 d^{\frac{2m}{d}-1} g_i y^{-i} \right| < 1.$$

Egész együtthatós polinom!





# A bizonyítás lépései

$$Q(x, y) := \sum_{i=0}^{\frac{n}{d}} d^{\frac{2m}{d}-1} f_{-i} x^i \pm \sum_{i=0}^{\frac{m}{d}} d^{\frac{2m}{d}-1} g_{-i} y^i = 0.$$

Megoldás  $x$ -re:  $\text{Res}_Y(F(X) - G(Y), Q(X, Y)) = 0$ .

Megoldás  $y$ -ra:  $\text{Res}_X(F(X) - G(Y), Q(X, Y)) = 0$ . Ahol

$$\text{Res}_Y(A(X, Y), B(X, Y)) = \begin{vmatrix} a_0(X) & \dots & \dots & a_r(X) \\ & \ddots & & \ddots \\ b_0(X) & \dots & a_0(X) & \dots & \dots & a_r(X) \\ & & & \ddots & & \ddots \\ & & & & \ddots & \\ b_0(X) & \dots & b_t(X) & & & b_r(X) \end{vmatrix}$$





# Algoritmus

Legyen  $t \in \mathbb{R}_{>0}$ . Tegyük fel, hogy  $p$  páratlan.

$$u_1(X) = \sum_{i=0}^{\frac{n}{p}} f_{-i} X^i \quad v_1(Y) = \sum_{i=0}^{\frac{m}{p}} g_{-i} Y^i$$

Ekkor

$$(u_1(x) - t)^p < F(x) < (u_1(x) + t)^p \text{ ha } x \notin [x_t^-, x_t^+],$$

$$(v_1(y) - t)^p < G(y) < (v_1(y) + t)^p \text{ ha } y \notin [y_t^-, y_t^+],$$

ahol

$$x_t^- = \min \left\{ \{0\} \cup \{x \in \mathbb{R} : F(x) - (u_1(x) - t)^p = 0 \text{ vagy } F(x) - (u_1(x) + t)^p = 0\} \right\},$$

$$x_t^+ = \max \left\{ \{0\} \cup \{x \in \mathbb{R} : F(x) - (u_1(x) - t)^p = 0 \text{ vagy } F(x) - (u_1(x) + t)^p = 0\} \right\},$$

$$y_t^- = \min \left\{ \{0\} \cup \{x \in \mathbb{R} : G(x) - (v_1(x) - t)^p = 0 \text{ vagy } G(x) - (v_1(x) + t)^p = 0\} \right\},$$

$$y_t^+ = \max \left\{ \{0\} \cup \{x \in \mathbb{R} : G(x) - (v_1(x) - t)^p = 0 \text{ vagy } G(x) - (v_1(x) + t)^p = 0\} \right\}.$$





# Algoritmus

Kapjuk, hogy

$$u_1(x) - t < F(x)^{1/p} < u_1(x) + t \text{ ha } x \notin [x_t^-, x_t^+],$$
$$v_1(y) - t < G(y)^{1/p} < v_1(y) + t \text{ ha } y \notin [y_t^-, y_t^+],$$

azaz

$$|u_1(x) - v_1(y)| < 2t.$$

$\Rightarrow x$  gyöke a

$$\text{Res}_Y(F(X) - G(Y), u(X) - v(Y) - T)$$

polinomnak, ahol  $-2t < T < 2t$  egy racionális szám, melynek nevezője a lemma alapján egy véges halmazból kerül ki.





# Algoritmus

A következő polinomok egész gyökeinek meghatározása adja az eredeti probléma megoldásait:

$$F(x) = G(k) \text{ ahol } k \in [y_t^-, y_t^+],$$

$$G(y) = F(k) \text{ valamelyen } k \in [x_t^-, x_t^+],$$

$$\text{Res}_Y(F(X) - G(Y), u(X) - v(Y) - T) = 0 \text{ ahol } T \in \mathbb{Q}, |T| < 2t$$

racionális szám, melynek nevezője osztja  $p^{\frac{2m}{p}-1}$ -t.

$t$  megválasztása: "baby-step giant-step", úgy, hogy az egyenletek száma alacsony legyen.





## Alkalmazás

$$x^2 - 3x + 5 = y^8 - y^7 + 9y^6 - 7y^5 + 4y^4 - y^3$$

Az  $u_1$  és  $v_1$  polinomok:

$$u_1(X) = X - \frac{3}{2},$$

$$v_1(Y) = Y^4 - \frac{1}{2}Y^3 + \frac{35}{8}Y^2 - \frac{21}{16}Y - \frac{1053}{128}.$$

$t = 1/16 \Rightarrow$  egyenletek száma= 158

Megoldások:

$$\{(-657, 5), (-3, -1), (0, 1), (3, 1), (6, -1), (660, 5)\}.$$





---

## Runge's condition

Consider a polynomial

$$P(X, Y) = \sum_{i=0}^m \sum_{j=0}^n a_{i,j} X^i Y^j,$$

where  $a_{i,j} \in \mathbb{Z}$  and  $m > 0, n > 0$ , which is irreducible in  $\mathbb{Q}[X, Y]$ .



## Runge's condition

Consider a polynomial

$$P(X, Y) = \sum_{i=0}^m \sum_{j=0}^n a_{i,j} X^i Y^j,$$

where  $a_{i,j} \in \mathbb{Z}$  and  $m > 0, n > 0$ , which is irreducible in  $\mathbb{Q}[X, Y]$ .  
In 1887 Runge gave an effective finiteness result concerning a class of Diophantine equation.

$$P(X, Y) = 0.$$





# Runge's condition

## Runge's condition

$P$  satisfies Runge's condition if at least one of the following conditions does not hold:

- $a_{i,n} = a_{m,j} = 0$  for all non-zero  $i$  and  $j$ ,
- for every term  $a_{i,j}X^iY^j$  of  $P$  one has  $ni + mj \leq mn$ ,
- the sum of all monomials  $a_{i,j}X^iY^j$  of  $P$  for which  $ni + mj = mn$  is up to a constant factor a power of an irreducible polynomial in  $\mathbb{Q}[X, Y]$ .





---

## Runge's condition

Let  $\lambda > 0$ .

- $P_\lambda(X, Y)$  : the sum of all terms  $a_{i,j}X^iY^j$  of  $P$  for which  $i + \lambda j$  is maximal,
- $\tilde{P}(X, Y)$  : the sum of all monomials of  $P$  which appear in any  $P_\lambda$  as  $\lambda$  varies.

$P$  satisfies Runge's condition unless there exists a  $\lambda$  so that  $\tilde{P} = P_\lambda$  is a constant multiple of a power of an irreducible polynomial.



## Runge's condition

Let  $\lambda > 0$ .

- $P_\lambda(X, Y)$  : the sum of all terms  $a_{i,j}X^iY^j$  of  $P$  for which  $i + \lambda j$  is maximal,
- $\tilde{P}(X, Y)$  : the sum of all monomials of  $P$  which appear in any  $P_\lambda$  as  $\lambda$  varies.

$P$  satisfies Runge's condition unless there exists a  $\lambda$  so that  $\tilde{P} = P_\lambda$  is a constant multiple of a power of an irreducible polynomial.

### Theorem (Runge (1887))

If  $P$  satisfies Runge's condition, then the Diophantine equation  $P(x, y) = 0$  has only a finite number of integer solutions.





## Example

$$P(X, Y) = X^2 - Y^8 - Y^7 - Y^2 - 3Y + 5,$$

- $P_\lambda(X, Y) = X^2$  if  $\lambda < \frac{1}{4}$ ,
- $P_\lambda(X, Y) = X^2 - Y^8$  if  $\lambda = \frac{1}{4}$ ,
- $P_\lambda(X, Y) = Y^8$  if  $\lambda > \frac{1}{4}$ ,

hence  $\tilde{P}(X, Y) = X^2 - Y^8 = (X - Y^4)(X + Y^4)$ .





---

## Related results

Schinzel (1969)

Masser (1980)

Hilliker and Straus (1983)

Ayad (1991)

Grytczuk and Schinzel (1991)

Walsh (1992)

Poulakis (1999)

Szalay (2000, 2002)

Tengely (2003)

Laurent and Poulakis (2004)

Levin (2004)

Beukers and Tengely (2005)

Sankaranarayanan and Saradha (2008)





# Bound for the solutions

## Theorem (Sz. T.)

Let  $F, G \in \mathbb{Z}[X]$  be monic polynomials with  $\deg F = n \leq \deg G = m$ , such that  $F(X) - G(Y)$  is irreducible in  $\mathbb{Q}[X, Y]$  and  $\gcd(n, m) > 1$ . Let  $d > 1$  be a divisor of  $\gcd(n, m)$ . If  $(x, y) \in \mathbb{Z}^2$  is a solution of the Diophantine equation  $F(x) = G(y)$ , then

$$\max\{|x|, |y|\} \leq d^{\frac{2m^2}{d}-m}(m+1)^{\frac{3m}{2d}}\left(\frac{m}{d}+1\right)^{\frac{3m}{2}}(h+1)^{\frac{m^2+mn+m}{d}+2m},$$

where  $h = \max\{H(F), H(G)\}$  and  $H(\cdot)$  denotes the classical height, that is the maximal absolute value of the coefficients.





## Sketch of the proof

### Lemma (Walsh)

There exist Puiseux expansions (in this case even Laurent expansions)

$$u(X) = \sum_{i=-\frac{n}{d}}^{\infty} f_i X^{-i} \text{ and } v(X) = \sum_{i=-\frac{m}{d}}^{\infty} g_i X^{-i}$$

of the algebraic functions  $U, V$  defined by  $U^d = F(X), V^d = G(X)$ , such that  $d^{2(n/d+i)-1} f_i \in \mathbb{Z}$  for all  $i > -\frac{n}{d}$ , similarly  $d^{2(m/d+i)-1} g_i \in \mathbb{Z}$  for all  $i > -\frac{m}{d}$ , and  $f_{-\frac{n}{d}} = g_{-\frac{m}{d}} = 1$ . Furthermore  $|f_i| \leq (H(F) + 1)^{\frac{n}{d}+i+1}$  for  $i \geq -\frac{n}{d}$  and  $|g_i| \leq (H(G) + 1)^{\frac{m}{d}+i+1}$  for  $i \geq -\frac{m}{d}$ .





---

## Sketch of the proof

$$F(X) = \left( \sum_{i=-\frac{n}{d}}^{\infty} f_i X^{-i} \right)^d, \quad G(Y) = \left( \sum_{i=-\frac{m}{d}}^{\infty} g_i Y^{-i} \right)^d$$

If  $|t|$  is large enough we obtain

$$\left| \sum_{i=1}^{\infty} d^{\frac{2m}{d}-1} f_i t^{-i} \right| < \frac{1}{2}$$

and

$$\left| \sum_{i=1}^{\infty} d^{\frac{2m}{d}-1} g_i t^{-i} \right| < \frac{1}{2}.$$





---

## Sketch of the proof

We have that  $F(x) = G(y)$ , so  $u(x)^d - v(y)^d = 0$ . Therefore we obtain that

$$(u(x) - v(y)) (u(x)^{d-1} + u(x)^{d-2}v(y) + \dots + v(y)^{d-1}) = 0,$$

if  $d$  is odd,

$$(u(x)^2 - v(y)^2) (u(x)^{d-2} + u(x)^{d-4}v(y)^2 + \dots + v(y)^{d-2}) = 0,$$

if  $d$  is even.



---

## Sketch of the proof

We have that  $F(x) = G(y)$ , so  $u(x)^d - v(y)^d = 0$ . Therefore we obtain that

$$(u(x) - v(y)) (u(x)^{d-1} + u(x)^{d-2}v(y) + \dots + v(y)^{d-1}) = 0,$$

if  $d$  is odd,

$$(u(x)^2 - v(y)^2) (u(x)^{d-2} + u(x)^{d-4}v(y)^2 + \dots + v(y)^{d-2}) = 0,$$

if  $d$  is even.

It follows that

$$u(x) = v(y) \text{ if } d \text{ is odd,}$$

$$u(x) = \pm v(y) \text{ if } d \text{ is even.}$$





---

## Sketch of the proof

$$0 = |u(x) \pm v(y)| = \left| \sum_{i=-\frac{n}{d}}^{\infty} f_i x^{-i} \pm \sum_{i=-\frac{m}{d}}^{\infty} g_i y^{-i} \right|.$$

If  $|x|$  and  $|y|$  are large enough we have

$$\left| \sum_{i=-\frac{n}{d}}^0 d^{\frac{2m}{d}-1} f_i x^{-i} \pm \sum_{i=-\frac{m}{d}}^0 d^{\frac{2n}{d}-1} g_i y^{-i} \right| < 1.$$

A polynomial having integer coefficients!





---

## Sketch of the proof

$$Q(x, y) := \sum_{i=0}^{\frac{n}{d}} d^{\frac{2m}{d}-1} f_{-i} x^i \pm \sum_{i=0}^{\frac{m}{d}} d^{\frac{2m}{d}-1} g_{-i} y^i = 0.$$

Solutions for  $x$ :  $\text{Res}_Y(F(X) - G(Y), Q(X, Y)) = 0$ .

Solutions for  $y$ :  $\text{Res}_X(F(X) - G(Y), Q(X, Y)) = 0$ . Where

$$\text{Res}_Y(A(X, Y), B(X, Y)) = \begin{vmatrix} a_0(X) & \dots & \dots & a_r(X) \\ & \ddots & & \ddots \\ b_0(X) & \dots & a_0(X) & \dots & \dots & a_r(X) \\ & \ddots & & \ddots & & \ddots \\ & & & & \ddots & \\ b_0(X) & \dots & b_t(X) & & & b_r(X) \end{vmatrix}$$





# Algorithm

Let  $t \in \mathbb{R}_{>0}$ . Suppose that  $p$  is odd.

$$u_1(X) = \sum_{i=0}^{\frac{n}{p}} f_{-i} X^i \quad v_1(Y) = \sum_{i=0}^{\frac{m}{p}} g_{-i} Y^i$$

Then we have

$$(u_1(x) - t)^p < F(x) < (u_1(x) + t)^p \text{ for } x \notin [x_t^-, x_t^+],$$

$$(v_1(y) - t)^p < G(y) < (v_1(y) + t)^p \text{ for } y \notin [y_t^-, y_t^+],$$

where

$$x_t^- = \min \{ \{0\} \cup \{x \in \mathbb{R} : F(x) - (u_1(x) - t)^p = 0 \text{ or } F(x) - (u_1(x) + t)^p = 0\} \},$$

$$x_t^+ = \max \{ \{0\} \cup \{x \in \mathbb{R} : F(x) - (u_1(x) - t)^p = 0 \text{ or } F(x) - (u_1(x) + t)^p = 0\} \},$$

$$y_t^- = \min \{ \{0\} \cup \{y \in \mathbb{R} : G(y) - (v_1(y) - t)^p = 0 \text{ or } G(y) - (v_1(y) + t)^p = 0\} \},$$

$$y_t^+ = \max \{ \{0\} \cup \{y \in \mathbb{R} : G(y) - (v_1(y) - t)^p = 0 \text{ or } G(y) - (v_1(y) + t)^p = 0\} \}.$$





---

# Algorithm

We obtain that

$$u_1(x) - t < F(x)^{1/p} < u_1(x) + t \text{ for } x \notin [x_t^-, x_t^+],$$
$$v_1(y) - t < G(y)^{1/p} < v_1(y) + t \text{ for } y \notin [y_t^-, y_t^+],$$

that is

$$|u_1(x) - v_1(y)| < 2t.$$

$\Rightarrow x$  is a root of the polynomial

$$\text{Res}_Y(F(X) - G(Y), u(X) - v(Y) - T),$$

for some rational number  $-2t < T < 2t$  with denominator dividing  $p^{\frac{2m}{p}-1}$ .





# Algorithm

It is sufficient to find all integral solutions of the following equations

$$F(x) = G(k) \text{ for some } k \in [y_t^-, y_t^+],$$

$$G(y) = F(k) \text{ for some } k \in [x_t^-, x_t^+],$$

$$\text{Res}_Y(F(X) - G(Y), u(X) - v(Y) - T) = 0 \text{ for some } T \in \mathbb{Q}, |T| < 2t$$

with denominator dividing  $p^{\frac{2m}{p}-1}$ .

The number of equations to be solved depends on  $t$ , a good choice can reduce the time of the computation.

We use the so-called "baby-step giant-step" algorithm to fix  $t$ .





---

## An application

$$x^2 - 3x + 5 = y^8 - y^7 + 9y^6 - 7y^5 + 4y^4 - y^3$$

We have

$$u(X) = X - \frac{3}{2},$$

$$v(Y) = Y^4 - \frac{1}{2}Y^3 + \frac{35}{8}Y^2 - \frac{21}{16}Y - \frac{1053}{128}.$$

$$t = 1/16 \Rightarrow \#\text{equations} = 158$$

Solutions:

$$\{(-657, 5), (-3, -1), (0, 1), (3, 1), (6, -1), (660, 5)\}.$$

